### <span id="page-0-0"></span>The Graph Abstract Data Type

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## **Outline**



### 1 [Introduction](#page-2-0)

- [Motivating Examples](#page-2-0)
- **•** [Graphs](#page-7-0)





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## <span id="page-2-0"></span>**Outline**



### 1 [Introduction](#page-2-0) **• [Motivating Examples](#page-2-0)**

[Graphs](#page-7-0)  $\bullet$ 

**[Graph Representations](#page-34-0)** 



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## <span id="page-3-0"></span>Königsberg Bridge Problem

### Question

Can we walk across all the bridges exactly once in returning back to the starting land area?



Recall **graphs** from the Discrete Mathematics course.

- Land *→* vertex
- Bridge *7→* edge



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## <span id="page-4-0"></span>Königsberg Bridge Problem

### Question (Eulerian Walk)

Can we walk across all the bridges exactly once in returning back to the starting land area?



Recall **graphs** from the Discrete Mathematics course.

- Land *→* vertex
- Bridge *7→* edge



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## <span id="page-5-0"></span>Eulerian Walk

### Euler's Theorem

*A connected graph has an Euler cycle if and only if every vertex has even degree.*

- **degree** of vertex *v*: number of neighbors of *v* in the graph.
- **connected**: there is a path connecting every two vertices in the graph.



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## <span id="page-6-0"></span>Eulerian Walk

### Euler's Theorem

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- **degree** of vertex *v*: number of neighbors of *v* in the graph.
- **connected**: there is a path connecting every two vertices in the graph.
- So, what about the answer to the Königberg Bridge Problem?



## <span id="page-7-0"></span>Definition of a Graph

### Graph

A graph  $G = (V, E)$  consists of two sets V and E, such that

- *V*: a finite, nonempty set of vertices;
- **•** *E*: a set of vertex pairs which are called edges.



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- **Undirected** graph: (*u, v*) and (*v, u*) represent the same edge.
- **Directed** graph: a directed vertex pair *< u, v >* has *u* as the tail and *v* as the head.



## <span id="page-9-0"></span>Definition of a Graph

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- **Directed** graph: a directed vertex pair *< u, v >* has *u* as the tail and *v* as the head.
	- $\bullet$   $\lt u, \nu$   $>$  and  $\lt v, u$   $>$  indicate different edges.

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<span id="page-10-0"></span>**Examples** 



 $V(G_1) = \{0, 1, 2, 3\}$  $E(G_1)=\{(0, 1), (0, 2),$  $(0, 3), (1, 2),$  $(1, 3), (2, 3)$ 

 $V(G_2) = \{0, 1, 2, 3, 4, 5, 6\}$  $E(G_2) = \{(0, 1), (0, 2), (1, 3),$  $(1, 4), (2, 5), (2, 6)$ 

 $V(G_3) = \{0, 1, 2\}$  $E(G_3) = \{ \le 0, 1 \ge 1, 0 \ge 1, 0 \le 1 \}$  $<1, 2>$ 



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## <span id="page-11-0"></span>Self-Loop & Multigrpah



**Graph with** a self-loop



multigraph

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- $\bullet$   $(v, v)$ : self-loop.
- multigraph: a graph with multiple occurrence of some edges.



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## <span id="page-12-0"></span>Complete Graph

### Complete Graph

An undirected graph  $G = (V, E)$  is a complete graph if any pair of vertices (*u, v*) is an edge in *E*.



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## <span id="page-13-0"></span>Complete Graph

### Complete Graph

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The number of edges in a complete undirected graph of *n* vertices:



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## <span id="page-14-0"></span>Complete Graph

### Complete Graph

An undirected graph  $G = (V, E)$  is a complete graph if any pair of vertices (*u, v*) is an edge in *E*.

- The number of edges in a complete undirected graph of *n* vertices:  $n(n-1)/2$ .
- The number of edges in a complete directed graph of *n* vertices:



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## <span id="page-15-0"></span>Complete Graph

### Complete Graph

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- The number of edges in a complete undirected graph of *n* vertices:  $n(n-1)/2$ .
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## <span id="page-16-0"></span>Subgraph and Induced Subgraph

 $\bullet$  If  $(u, v)$  is an edge in  $E(G)$ , then the vertices *u* and *v* are adjacent and that the edge  $(u, v)$  is incident on vertices  $u$  and  $v$ .

### Subgraph

A subgraph of *G* is a graph *G*<sup>*'*</sup> such that  $V(G) \subseteq V(G)$  and  $E(G') \subseteq EG$ .

### Induced Subgraph

A graph *G ′* is an induced subgraph of *G* if *G ′* is a subgraph of *G* and for any two vertices  $u, v \in V(G')$ ,  $(u, v) \in E(G)$  if and only if  $(u, v) \in E(G')$ .



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## **Examples**





<span id="page-18-0"></span>Examples



 $\bullet$  *H*<sub>1</sub>, *H*<sub>2</sub>: subgraphs of *G*.



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- $\bullet$  *H*<sub>1</sub>*, H*<sub>2</sub>: subgraphs of *G*.
- $H_1$  is an induced subgraph of *G*, but  $H_2$  is NOT.



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## <span id="page-20-0"></span>Path  $(1/2)$

### Path

A (directed or undirected) path from vertex *u* to vertex *v* in graph *G* is a sequence of vertices u,  $i_1, i_2, \ldots, i_k, v$ , such that  $(u, i_1), (i_1, i_2), \ldots, (i_k, v)$ are edges in *E*(*G*).



## <span id="page-21-0"></span>Path  $(1/2)$

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## <span id="page-22-0"></span>Path  $(1/2)$

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## Path (2/2)

- The length of a path is the number of edges on it.
- A simple path a path in which all vertices, except possibly the first and the last, are distinct.



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## Path (2/2)

- The length of a path is the number of edges on it.
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- A cycle is a simple path in which the first and last vertices are the same.
	- A simple path from *v* to *v*.



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## <span id="page-26-0"></span>Connected and Connected Component

### **Connected**

- In an undirected graph *G*, two vertices *u* and *v* are connected iff there is a path in *G* from *u* to *v*.
- An undirected graph is **connected** iff for every pair of district vertices *u* and *v* in *V*(*G*) there is a path from *u* to *v* in *G*.

### Connected Component

A **connected component** (or simply a component) *H* of an undirected graph is a maximal connected subgraph.



## <span id="page-27-0"></span>Connected and Connected Component

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A **connected component** (or simply a component) *H* of an undirected graph is a maximal connected subgraph.

#### tree

A tree is a connected acyclic (i.e., has no cycles) graph.

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## <span id="page-28-0"></span>Example of Connected Components



## <span id="page-29-0"></span>Strongly Connected Graph

### Strongly Connected Graph

A directed graph G is said to be **strongly connected** iff for every pair of district vertices  $u, v \in V(G)$ , there is directed path from *u* to *v* and also from *v* to *u*.



## <span id="page-30-0"></span>Strongly Connected Graph

### Strongly Connected Graph

A directed graph G is said to be **strongly connected** iff for every pair of district vertices  $u, v \in V(G)$ , there is directed path from *u* to *v* and also from *v* to *u*.

### Strongly Connected Component

A strongly connected component is a maximal subgraph that is strongly connected.



## <span id="page-31-0"></span>Strongly Connected Components





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### Vertex Degree

- The degree of a vertex is the number of edges **incident to** that vertex.
- For a directed graph *G*,
	- The in-degree of a vertex is the number of edges for which vertex is head.
	- the out-degree of a vertex is the number of edges for which the vertex is the tail.



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Let *d<sup>i</sup>* be the degree of vertex *i* in an *n*-vertex graph  $G = (V, E)$ , then

$$
|E|=\frac{1}{2}\sum_{i=1}^n d_i.
$$

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## <span id="page-34-0"></span>**Outline**







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### <span id="page-35-0"></span>Graph Representations

Two most commonly used representation for a graph:



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## <span id="page-36-0"></span>Graph Representations

Two most commonly used representation for a graph:

- Adjacency Matrices
- Adjacency Lists



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## <span id="page-37-0"></span>Graph Representations

Two most commonly used representation for a graph:

- **Adjacency Matrices**
- **Adjacency Lists**

#### • The choice of the representation:

- the application
- the functions one expects to perform on the graph
- characteristics of the input graph



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## <span id="page-38-0"></span>Adjacency Matrix

The adjacency matrix of an *n*-vertex graph *G* is a two-dimensional *n × n* array a, with the property that

• 
$$
a[i][j] = 1
$$
 iff  $(i, j) \in E(G)$ ;

•  $a[i][j] = 0$  iff there is no such edge  $(i, j)$  in *G*.

#### Remark

The adjacency matrix for an undirected graph is symmetric.

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## <span id="page-39-0"></span>Adjacency Matrix (2/2)

- For an undirected graph the degree of any vertex *i* is its row sum.
- For an directed graph the row sum is its out-degree and the column sum is its in-degree.



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## <span id="page-40-0"></span>Adjacency Matrices (Examples)

The adjacency matrix of  $G_1$ 





The adjacency matrix of  $G_4$ 



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## <span id="page-41-0"></span>Adjacency Lists

The *n* rows of the adjacency matrix are represented as *n* chains.



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## <span id="page-42-0"></span>Adjacency Lists

The *n* rows of the adjacency matrix are represented as *n* chains.

One chain for each vertex in *G*.



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## <span id="page-43-0"></span>Adjacency Lists

- The *n* rows of the adjacency matrix are represented as *n* chains.
- One chain for each vertex in *G*.
- The nodes in chain *i* represent the vertices that are adjacent from vertex *i*.



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## <span id="page-44-0"></span>Adjacency Lists

- The *n* rows of the adjacency matrix are represented as *n* chains.
- One chain for each vertex in *G*.
- The nodes in chain *i* represent the vertices that are adjacent from vertex *i*.
- The data field of a chain node stores the index of an adjacent vertex.

## <span id="page-45-0"></span>Adjacency Lists Examples



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## <span id="page-46-0"></span>Remark: Weighted Edges

• In many applications, the edges of a graph have weights associated with them.



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## <span id="page-47-0"></span>Remark: Weighted Edges

- In many applications, the edges of a graph have weights associated with them.
	- importance, costs, distance, etc.
- The adjacency matrix entries a[i][j] would keep this information.
- When adjacency lists are used, we can introduce an additional field **weight** in the list nodes.



# <span id="page-48-0"></span>**Discussions**



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