

# Randomized Algorithms

## Discrete Random Variables and Expectation

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# Review

- Why randomized algorithms?
- Types of randomized algorithms.

# Randomized algorithms

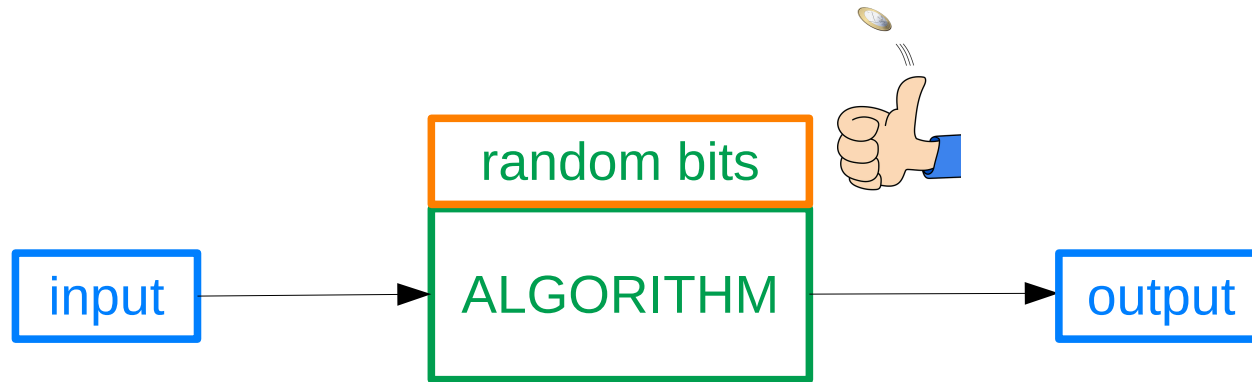


figure freely downloaded from <https://pixabay.com/illustrations/coin-flipping-coin-hand-flip-flick-5822271/>

*Randomized Algorithms, CSIE, TKU, Taiwan*

# Why?

- Randomized algorithms are
  - often much *simpler* than the best known deterministic ones.
  - often much *more efficient* (faster or using less space) than the best known deterministic ones.

# Two types of randomized algorithms

- The **accuracy** is guaranteed.
  - Las Vegas algorithms.
- The **running time** is guaranteed.
  - Monte Carlo algorithms.

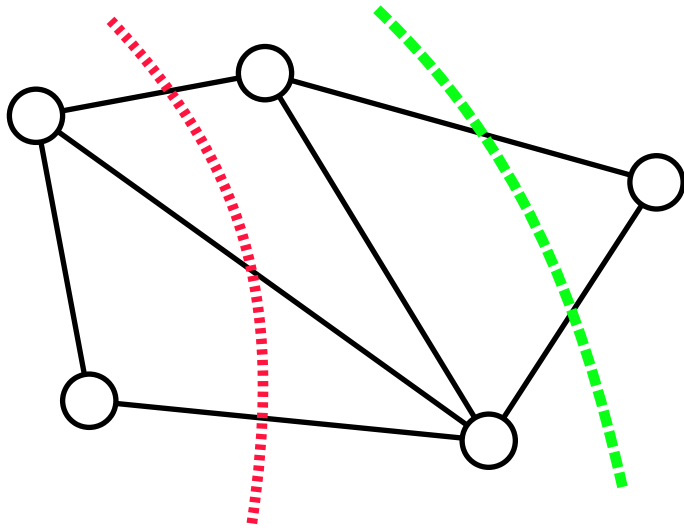


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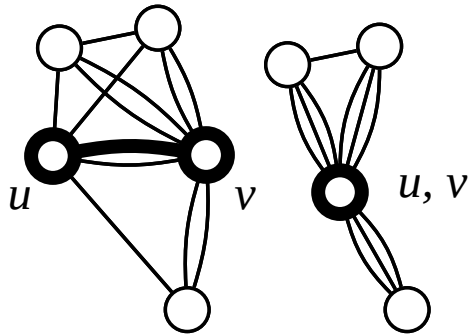
# Min-Cut



$$|V| = n, |E| = m$$

- A graph  $G = (V, E)$  and its two “cuts”.
  - **Cut**: a partition of the vertices in  $V$  into two non-empty, disjoint sets  $S$  and  $T$  such that
    - $S \cup T = V$
- The **cutset** of a cut:
  - $\{uv \in E \mid u \in S, v \in T\}$ .
- The size of the cut:
  - the cardinality of its cutset.

# Edge contraction



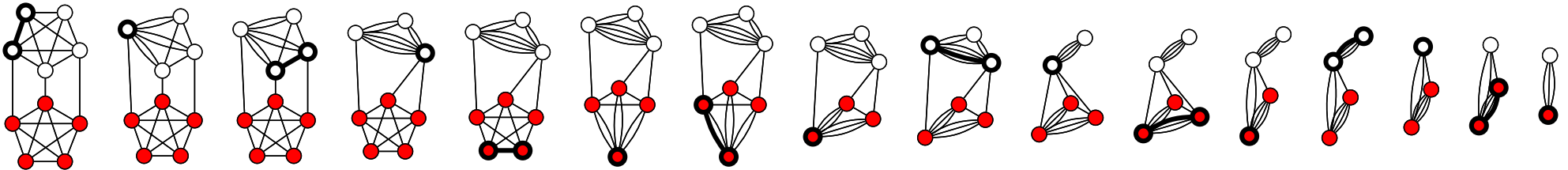
$$e = (u, v)$$

$$G \rightarrow G/e$$

# Karger's edge-contraction algorithm (1993)

```
Procedure contract ( $G = (V, E)$ ):  
while  $|V| > 2$ :  
    choose  $e \in E$  uniformly at random  
     $G \leftarrow G/e$   
return the only cut in  $G$ 
```

Time complexity:  $O(m)$  or  $O(n^2)$ .



By Thore Husfeldt - Created in python using the networkx library for graph manipulation, neato for layout, and TikZ for drawing.,  
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# Analysis

- $C$ : a specific cut of  $G$ .
- $k$ : the size of the cut  $C$ .
- The minimum degree of  $G$  must be  $\geq k$ . (WHY?)
  - So,  $|E| \geq nk/2$ .
- The probability that the algorithm picks an edge from  $C$  to contract is

$$\frac{k}{|E|} \leq \frac{k}{nk/2} = \frac{2}{n}.$$

# Analysis (contd.)

- Let  $p_n$  be the probability that the algorithm on an  $n$ -vertex graph avoids  $C$ .

- Then,

$$p_n \geq \left(1 - \frac{2}{n}\right) \cdot p_{n-1}$$

- The recurrence can be expanded as

$$p_n \geq \prod_{i=0}^{n-3} \left(1 - \frac{2}{n-i}\right) = \frac{n-2}{n} \cdot \frac{n-3}{n-1} \cdots \frac{2}{4} \cdot \frac{1}{3} = \frac{1}{\binom{n}{2}}.$$

# Analysis (contd.)

- Repeat the contract algorithm for  $T = \binom{n}{2} \ln n$  times, and then choose the minimum of them.
- The probability of NOT finding a min-cut is  $\left[1 - \binom{n}{2}^{-1}\right]^T \leq \frac{1}{e^{\ln n}} = \frac{1}{n}$ .
- Let's take a look at the exponential function  $e^x$ .

# Facts on $e^x$

$$e^x := \sum_{k=0}^{\infty} \frac{x^k}{k!} = 1 + x + \frac{x^2}{2} + \frac{x^3}{6} + \dots$$

$$\therefore e = \sum_{k=0}^{\infty} \frac{1}{k!}$$

$$\left(1 + \frac{1}{n}\right)^n = \binom{n}{0} 1 + \binom{n}{1} \frac{1}{n} + \binom{n}{2} \frac{1}{n^2} + \binom{n}{3} \frac{1}{n^3} + \dots + \binom{n}{n} \frac{1}{n^n}$$

$$\lim_{n \rightarrow \infty} \left(1 + \frac{1}{n}\right)^n = e$$

# Probability basics

A *random variable*  $X$  on a sample space  $\Omega$  is a real-valued function. That is,

$$X : \Omega \mapsto \mathbf{R}$$

A *discrete random variable* is a random variable that takes on only a finite or countably infinite number of values.

# Kolmogorov axioms

- 1. For any event  $A \subset S$ ,  $P(A) \geq 0$ .
- 2.  $\Pr[S] = 1$ .
- 3. If  $A_1, A_2, \dots$  are mutually exclusive events, then
$$\Pr[A_1 \cup A_2 \cup \dots] = \Pr[A_1] + \Pr[A_2] + \dots$$

# Useful theorems

- $\Pr[\emptyset] = 0$  for any experiment.
- For any event  $A \subseteq S$ ,  $\Pr[A] = 1 - \Pr[\bar{A}]$ .
- If  $A \subseteq S$ ,  $B \subseteq S$  are any two events, then  
 $\Pr[A \cup B] = \Pr[A] + \Pr[B] - \Pr[A \cap B]$ .
- If  $A \subset B$ , then  $\Pr[A] \leq \Pr[B]$ .

# Conditional probability

- In an experiment with sample space  $S$ , let  $B$  be any event such that  $\Pr[B] > 0$ .
- Then the conditional probability of  $A$  occurring, given that  $B$  has occurred, is

$$\Pr[A \mid B] = \frac{\Pr[A \cap B]}{\Pr[B]}$$

for any  $A \subset S$ .



# Example

- Suppose you are going to buy milk in a supermarket.
  - There are a total of 40 boxes for you to choose from.
  - 10 of them are corrupted (not visible on the outside).
  - Then, you are asked to buy two boxes of milk.  
**What is the probability that both boxes are good?**

# Example (contd.)

- $A$ : the event that the first box you choose is good.  
 $B$ : the event that the second box you choose is good.

Then

$$\Pr[A] = \frac{30}{40}$$

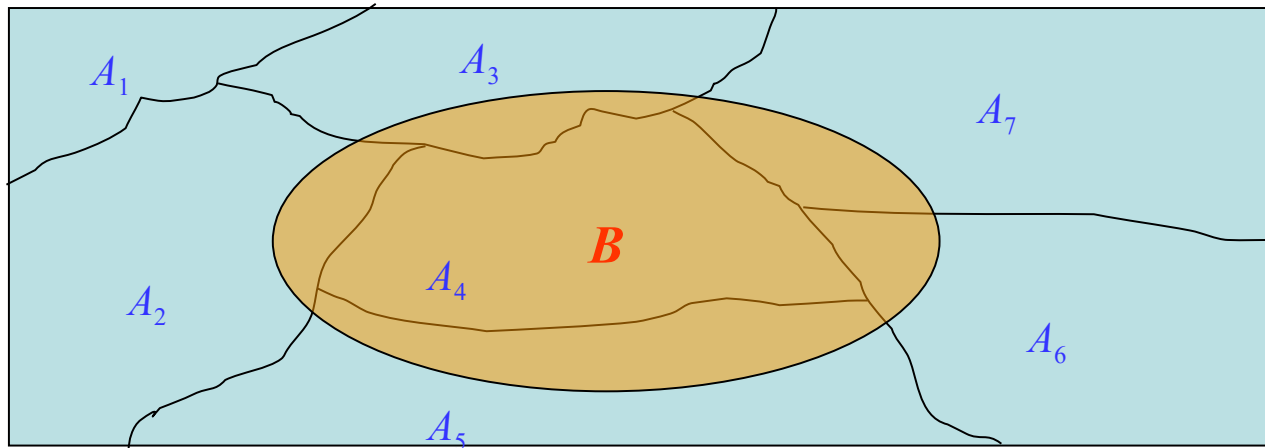
$$\Pr[B \mid A] = \frac{29}{39}$$

$$\Pr[A \cap B] = \Pr[A] \cdot \Pr[B \mid A] = \frac{30}{40} \cdot \frac{29}{39} = \frac{87}{156}.$$

# Theorem of total probability

- If  $A_1, A_2, \dots, A_n$  is a partition of  $S$ , and  $B$  is any event, then

$$\Pr[B] = \sum_{i=1}^n \Pr[B|A_i] \Pr[A_i].$$



# Bayes' theorem

- From the theorem of total probability, and granted that  $A_1, A_2, \dots, A_n$  is a partition of  $S$ , we have

$$\begin{aligned}\Pr[A_i | B] &= \frac{\Pr[A_i \cap B]}{\Pr[B]} \\ &= \frac{\Pr[A_i] \cdot \Pr[B | A_i]}{\sum_{i=1}^n \Pr[A_i] \cdot \Pr[B | A_i]}\end{aligned}$$

- This result is known as *Bayes' theorem*.

# Example

- Assuming a jury selected to participate in a criminal trial.
- Whether the defendant is guilty or not, there is a 95% chance of making the correct verdict.
- It is also assumed that the local police law enforcement is very strict, such that 99% of the people being tried are actually guilty.
- **If a jury is known to sentence a defendant not guilty, what is the probability that the defendant is really not guilty?**

# Example (contd.)

- $A_1$ : the defendant is guilty
- $A_2 = \bar{A}_1$  : the defendant is not guilty.
- Let  $B$  be the event that the defendant is sentenced to unguilty.
- We want to know  **$\Pr[A_2 | B]$** .

# Example (contd.)

$$\begin{aligned}\Pr[A_2|B] &= \frac{\Pr[A_2]\Pr[B|A_2]}{\Pr[A_1]\Pr[B|A_1] + \Pr[A_2]\Pr[B|A_2]} \\ &= \frac{(0.01)(0.95)}{(0.99)(0.05) + (0.01)(0.95)} \\ &= 0.161\end{aligned}$$

- Before the sentence, this defendant is supposed to be unguilty with probability 1%.
- After the sentence of unguilty, the probability is increased to be **16.1%**.

# Independent events

If  $A \subset S$  and  $B \subset S$  are any two events with nonzero probabilities,  $A$  and  $B$  are called independent if and only if  $\Pr[A \cap B] = \Pr[A] \cdot \Pr[B]$

That is,  $\Pr[A] = \Pr[A | B]$  and  $\Pr[B] = \Pr[B | A]$ .



# Independent trials

- An experiment is said to consist of  $n$  ***independent*** trials if and only if
  - $S = T_1 \times T_2 \times \cdots \times T_n$ .
  - For every  $(x_1, x_2, \dots, x_n) \in S$ ,  
$$\Pr[\{(x_1, x_2, \dots, x_n)\}] = \Pr[\{x_1\}] \cdot \Pr[\{x_2\}] \cdots \Pr[\{x_n\}],$$
where  $\Pr[\{x_i\}]$  is the probability of  $x_i \in T_i$  occurring on trial  $i$ .

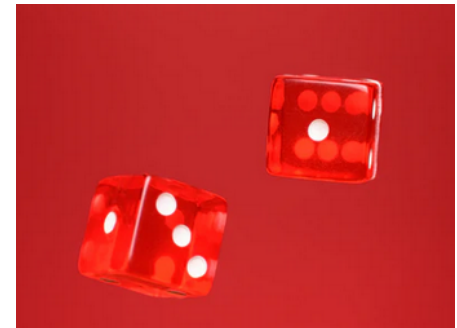
# Expectation

- The expectation of a discrete random variable  $X$ , denoted by  $\mathbf{E}[X]$ , is

$$\mathbf{E}[X] = \sum_i i \cdot \Pr[X = i]$$

- Example: Let  $X$  denote the sum of of dices:

$$\mathbf{E}[X] = \frac{1}{36} \cdot 2 + \frac{2}{36} \cdot 3 + \frac{3}{36} \cdot 4 + \cdots + \frac{1}{36} \cdot 12 = 7.$$



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# Linearity of Expectation

- For any finite collection of discrete random variables  $X_1, X_2, \dots, X_n$  with finite expectations,

$$\mathbf{E} \left[ \sum_{i=1}^n X_i \right] = \sum_{i=1}^n \mathbf{E}[X_i].$$

- For any constant  $c$  and discrete random variable  $X$ ,

$$\mathbf{E}[cX] = c \cdot \mathbf{E}[X].$$

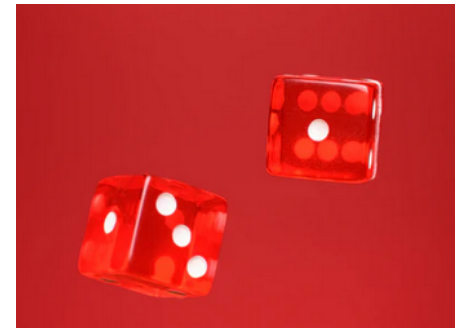
- Why is it useful?

# Example

- Consider the dice-throwing example again.
  - $X_1$  : the outcome of die 1
  - $X_2$  : the outcome of die 2

$$\mathbf{E}[X_1] = \mathbf{E}[X_2] = \frac{1}{6} \cdot \sum_{j=1}^6 j = \frac{7}{2}$$

$$\mathbf{E}[X] = \mathbf{E}[X_1 + X_2] = 7.$$



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# Bernoulli random variable

- Suppose we run an experiment that succeeds with probability  $p$  and fails with probability  $1-p$ .

$$Y = \begin{cases} 1 & \text{if the experiment succeeds,} \\ 0 & \text{otherwise.} \end{cases}$$

- $Y$ : Bernoulli random variable.
  - or *indicator random variable*.



# Binomial random variable

- A binomial random variable  $X$  with parameters  $n$  and  $p$ , denoted by  $B(n, p)$ , is defined as

$$\Pr[X = j] = \binom{n}{j} p^j (1 - p)^{n-j}.$$

for  $j = 0, 1, 2, \dots, n$ .

- Exercise: Show that  $\sum_{j=0}^n \Pr[X = j] = 1$ .

# Binomial random variable (contd.)

- $\mathbf{E}[X] = \sum_{j=0}^n j \binom{n}{j} p^j (1-p)^{n-j}$ 
$$= \sum_{j=0}^n j \frac{n!}{j!(n-j)!} p^j (1-p)^{n-j}$$
$$= \sum_{j=1}^n \frac{n!}{(j-1)!(n-j)!} p^j (1-p)^{n-j}$$
$$= np \sum_{j=1}^n \frac{(n-1)!}{(j-1)!((n-1)-(j-1))!} p^{j-1} (1-p)^{(n-1)-(j-1)}$$
$$= np \sum_{k=0}^{n-1} \frac{(n-1)!}{k!((n-1)-k)!} p^k (1-p)^{(n-1)-k}$$
$$= np \sum_{k=0}^{n-1} \binom{n-1}{k} p^k (1-p)^{(n-1)-k}$$
$$= np.$$

# Binomial random variable (contd.)

- $\mathbf{E}[X] = \sum_{j=0}^n j \binom{n}{j} p^j (1-p)^{n-j}$

$$= \sum_{j=0}^n j \frac{n!}{j!(n-j)!} p^j (1-p)^{n-j}$$


$$= \sum_{j=1}^n \frac{n!}{(j-1)!(n-j)!} p^j (1-p)^{n-j}$$

$$= np \sum_{j=1}^n \frac{(n-1)!}{(j-1)!((n-1)-(j-1))!} p^{j-1} (1-p)^{(n-1)-(j-1)}$$

$$= np \sum_{k=0}^{n-1} \frac{(n-1)!}{k!((n-1)-k)!} p^k (1-p)^{(n-1)-k}$$

$$= np \sum_{k=0}^{n-1} \binom{n-1}{k} p^k (1-p)^{(n-1)-k}$$

$$= np.$$


$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k}.$$



# Let's make it simpler!

- Denote a set of  $n$  Bernoulli random variables  $X_1, X_2, \dots, X_n$ .
  - $X_i = 1$  if the  $i$ th trial is successful and 0 otherwise.

$$\mathbf{E}[X] = \mathbf{E} \left[ \sum_{i=1}^n X_i \right] = \sum_{i=1}^n \mathbf{E}[X_i] = np.$$